

## B9145: Problem Set 2

Due: Oct 30, 11:59pm

Carefully follow submission instructions announced on Canvas.

**Question 2.1** (Minimax bounds for estimation (30 points)): We derive information theoretic lower bounds for statistical estimation problems, analogous to those for stochastic optimization we saw in class. For a class of distributions  $\mathcal{P}$ , let  $\theta : \mathcal{P} \rightarrow \mathbb{R}^d$  be the statistical functional of interest;  $\theta(P)$  is often called the “parameter”. Let  $d$  be a metric on  $\Theta := \{\theta(P) : P \in \mathcal{P}\}$ , and let  $\Phi : \mathbb{R}_+ \rightarrow \mathbb{R}_+$  be a non-decreasing function such that  $\Phi(0) = 0$ . For  $n$  observations  $X_i \stackrel{\text{iid}}{\sim} P$ , we measure performance of an estimator  $\hat{\theta}_n(X_1, \dots, X_n)$  by

$$\sup_{P \in \mathcal{P}} \mathbb{E}_{X_1^n \sim P} \left[ \Phi \left( d(\hat{\theta}_n(X_1^n), \theta(P)) \right) \right].$$

The minimax risk for estimation is given by

$$\mathfrak{M}_n(\mathcal{P}, \Phi \circ d) := \inf_{\hat{\theta}} \sup_{P \in \mathcal{P}} \mathbb{E}_P \left[ \Phi \left( d(\hat{\theta}_n(X_1^n), \theta(P)) \right) \right],$$

where the infimum is taken over all measurable functions of  $X_1, \dots, X_n$ .

For parts (a)-(b), you may give a concise derivation based on results from class.

(a) Derive Le Cam’s method: for any fixed  $\delta > 0$ , and  $P_1, P_{-1} \in \mathcal{P}$  such that  $d(\theta(P_1), \theta(P_{-1})) \geq 2\delta$ ,

$$\mathfrak{M}_n(\mathcal{P}, \Phi \circ d) \geq \frac{\Phi(\delta)}{2} (1 - \|P_1^n - P_{-1}^n\|_{\text{TV}}).$$

(b) Derive Assouad’s method. Let  $\mathcal{V} := \{-1, +1\}^d$  be the binary hypercube, and let  $\{P_v^n\}_{v \in \mathcal{V}}$  be a collection of distributions on  $X_1^n$ . We say  $\{P_v^n\}_{v \in \mathcal{V}}$  is  $\delta$ -separated in the Hamming distance if there exists a mapping  $\hat{v} : \Theta \mapsto \mathcal{V}$  such that

$$\Phi(d(\theta, \theta(P_v))) \geq \delta \sum_{j=1}^d \mathbf{1}\{\hat{v}(\theta)_j \neq v_j\}.$$

Define  $P_{+j}^n := \frac{1}{2^{d-1}} \sum_{v: v_j=1} P_v^n$  and  $P_{-j}^n := \frac{1}{2^{d-1}} \sum_{v: v_j=-1} P_v^n$ . Then, we have

$$\mathfrak{M}_n(\mathcal{P}, \Phi \circ d) \geq \frac{\delta}{2} \sum_{j=1}^d \left( 1 - \|P_{+j}^n - P_{-j}^n\|_{\text{TV}} \right)$$

whenever  $\{P_v^n\}_{v \in \mathcal{V}}$  is  $\delta$ -separated in Hamming distance.

- (c) Consider the normal location model  $\mathcal{P}_{\sigma^2} := \{N(\theta, \sigma^2 I) : \theta \in \mathbb{R}^d\}$ , where  $I$  is the  $d$ -by- $d$  dimensional identity matrix, and  $\sigma^2 > 0$  is a fixed variance. We're interested in estimating the location parameter  $\theta$  in the squared Euclidean distance  $\|\cdot\|_2^2$ . Show the following lower bound

$$\mathfrak{M}_n(\mathcal{P}_{\sigma^2}, \|\cdot\|_2^2) \geq \frac{d\sigma^2}{16n}. \quad (1)$$

- (d) Argue that the lower bound (1) is tight up to numerical constants.

**Question 2.2** (Differentially private estimation (50 points)): We study estimation under a privacy constraint, when the data collector cannot be trusted with sensitive information. Instead of observing true data  $X_i \in \mathcal{X}$ , a perturbed version  $Z_i \in \mathcal{Z}$  is viewed; given  $X = x$ , we write  $Z \sim Q(\cdot | X = x)$ , and call  $Q$  a “channel”. For  $\alpha > 0$ , we say  $Z_i$  is  $\alpha$ -differentially private if for any measurable subset  $A \subset \mathcal{Z}$  and any pair  $x, x' \in \mathcal{X}$ ,

$$\frac{Q(Z \in A | X = x)}{Q(Z \in A | X = x')} \leq \exp(\alpha). \quad (2)$$

Intuitively, differential privacy asks that  $x$  and  $x'$  are similarly likely to have generated the observed signal  $Z$ . Letting  $q(z | x) := Q(Z = z | X = x)$  be the conditional density of  $Z | X$ , the condition (2) is equivalent to  $\frac{q(z|x)}{q(z|x')} \leq e^\alpha$  for all  $x, x' \in \mathcal{X}$ , and almost surely all  $z \in \mathcal{Z}$ . In what follows, we assume  $\alpha < 1$ .

As we will show, differential privacy acts as a contraction on probabilities. For arbitrary probabilities  $P_1, P_2$  on  $\mathcal{X}$ , let densities  $p_1$  and  $p_2$  be their densities w.r.t. a base measure  $\mu$ ; you may treat this as a continuous density for convenience. Define the *marginal* distributions

$$M_i(Z \in A) := \int_{\mathcal{X}} Q(Z \in A | X = x) p_i(x) d\mu(x), \quad i \in \{1, 2\}.$$

We will prove there is a universal (numerical) constant  $C < \infty$  such that for any  $P_1, P_2$ ,

$$D_{\text{kl}}(M_1 \| M_2) + D_{\text{kl}}(M_2 \| M_1) \leq C(e^\alpha - 1)^2 \|P_1 - P_2\|_{\text{TV}}^2. \quad (3)$$

We show this result assuming  $\mathcal{Z} = \{1, \dots, k\}$  for some finite  $k \in \mathbb{N}$ ; this is without loss of generality, but you don't have to justify this.

- (a) Recall the definition of the total variation distance  $\|P_1 - P_2\|_{\text{TV}} = \sup_{A \subset \mathcal{X}} \{P_1(A) - P_2(A)\}$ . Show  $\|P_1 - P_2\|_{\text{TV}} = \frac{1}{2} \int |p_1(x) - p_2(x)| d\mu(x)$ .
- (b) Define  $m_j(z) := \int q(z | x) p_j(x) d\mu(x)$ , prove that for a universal constant  $c < \infty$ ,

$$|m_1(z) - m_2(z)| \leq c(e^\alpha - 1) \inf_{x \in \mathcal{X}} q(z | x) \cdot \|P_1 - P_2\|_{\text{TV}}.$$

- (c) Show the result (3) when  $\mathcal{Z} = \{1, \dots, k\}$  for some finite  $k \in \mathbb{N}$ .

**Hint** Use the following simple inequality: for any  $a, b > 0$ , we have  $|\log \frac{a}{b}| \leq \frac{|a-b|}{\min\{a, b\}}$ . To see this, use  $\log(1+x) \leq x$  to note

$$\log \frac{a}{b} = \log \left(1 + \frac{a}{b} - 1\right) \leq \frac{a-b}{b} \quad \text{and} \quad \log \frac{b}{a} \leq \frac{b-a}{a}.$$

We now use the inequality (3) to prove minimax lower bounds for differentially private estimation. Consider a survey data on individuals  $i = 1, \dots, n$ , where we ask each individual about illicit drug use:  $X_i = 1$  if person  $i$  uses illicit drugs, 0 otherwise ( $\mathcal{X} = \{0, 1\}$ ). Define  $\theta(P) = P(X = 1) = \mathbb{E}_P[X]$ . To protect privacy, we perturb each answer  $X_i$  in a  $\alpha$ -differentially private manner, and use  $Z_i$ 's as our data.

To make sure everyone feels suitably private, assume  $\alpha < 1/2$ ; in this case,  $(e^\alpha - 1)^2 \leq 2\alpha^2$ . Let  $\mathcal{Q}_\alpha$  be the family of all  $\alpha$ -differentially private channels, and let  $\mathcal{P}$  denote the Bernoulli distributions with parameter  $\theta(P) = P(X_i = 1) \in [0, 1]$ . We consider the minimax risk for private estimation of the proportion  $\theta(P)$

$$\mathfrak{M}_n(\theta(\mathcal{P}), |\cdot|, \alpha) := \inf_{Q \in \mathcal{Q}_\alpha} \inf_{\hat{\theta}} \sup_{P \in \mathcal{P}} \mathbb{E} \left[ |\hat{\theta}(Z_1, \dots, Z_n) - \theta(P)| \right],$$

where the infimum is over (differentially private) channels  $Q$  and estimators  $\hat{\theta}$ , and the expectation is taken with respect to both the  $X_i$  (according to  $P$ ) and the  $Z_i$  (according to  $Q(\cdot | X_i)$ ).

(d) Use Le Cam's method to argue that whenever  $P_1, P_2$  satisfy  $|\theta(P_1) - \theta(P_2)| \geq \delta$ ,

$$\mathfrak{M}_n(\theta(\mathcal{P}), |\cdot|, \alpha) := \inf_{Q \in \mathcal{Q}_\alpha} \inf_{\hat{\theta}} \sup_{P \in \mathcal{P}} \mathbb{E} \left[ |\hat{\theta}(Z_1, \dots, Z_n) - \theta(P)| \right] \geq \frac{\delta}{2} \inf_{Q \in \mathcal{Q}_\alpha} [1 - \|M_1^n - M_2^n\|_{\text{TV}}].$$

Then, use inequality (3) to show that for some universal constant  $c' > 0$ .

$$\mathfrak{M}_n(\theta(\mathcal{P}), |\cdot|, \alpha) \geq \frac{c'}{\sqrt{n\alpha^2}}.$$

(e) Give a rate-optimal estimator for this problem. i.e., define a  $\alpha$ -differentially private channel  $Q$  and an estimator  $\hat{\theta}$  such that  $\mathbb{E}[|\hat{\theta}(Z_1^n) - \theta|] \leq C'/\sqrt{n\alpha^2}$ , where  $C' > 0$  is a universal constant.

**Hint** Consider perturbing the data with probability  $1 - q_\alpha$ , where  $q_\alpha = e^\alpha / (1 + e^\alpha)$ . Note that  $(2q_\alpha - 1)^{-2} = \left(\frac{e^\alpha + 1}{e^\alpha - 1}\right)^2 \approx 4/\alpha^2$  for  $\alpha \approx 0$ .

**Question 2.3** (Adversarial robustness for linear logistic regression (10 points)): Consider a binary classification problem with label  $y \in \{-1, +1\}$  and features  $x \in \mathbb{R}^d$ . We study the logistic regression loss  $\ell(\theta; x, y) = -\log \sigma(y\theta^\top x)$ , where  $\sigma(a) = \frac{1}{1 + \exp(-a)}$ . Derive an alternative form for the adversarial loss:

$$\max_{\bar{x} \in \mathbb{R}^d: \|\bar{x} - x\|_\infty \leq \epsilon} \ell(\theta; \bar{x}, y) = -\log \sigma \left( y\theta^\top x - \epsilon \|\theta\|_1 \right).$$

Give an interpretation of this result.

**Question 2.4** (ICLR 2020 Vision talk (10 points)): Watch Ruha Benjamin's ICLR talk on "Reimagining the default settings of technology and society" via the url [https://iclr.cc/virtual\\_2020/speaker\\_3.html](https://iclr.cc/virtual_2020/speaker_3.html). In 2-3 sentences, discuss how this may relate to your research, or other professional activities.